



# Omid: A Transactional Framework for HBase

Francisco Perez-Sorrosal

Ohad Shacham

Hadoop Summit SJ

June 29th, 2016

# Outline



- Background
- Basic Concepts
- Use cases
- Architecture
- Transaction Management
- High Availability
- Performance
- Summary

# Background

- New Big data apps → new requirements:
  - Low-latency
  - Incremental data processing
  - e.g. Percolator
- Multiple clients updating same data concurrently
  - **Problem:** Conflicts/Inconsistencies may arise
  - **Solution: Transactional Access to Data**

# Background

- Transaction → Abstract UoW to manage data with certain guarantees
  - ACID
  - Relational databases
- Big data → NoSQL datastores → Transactions in NoSQL
  - Hard to Scale
    - Data partition
    - Data replication
  - Relaxed Guarantees:
    - e.g. Atomicity, Consistency

# Omid is a...

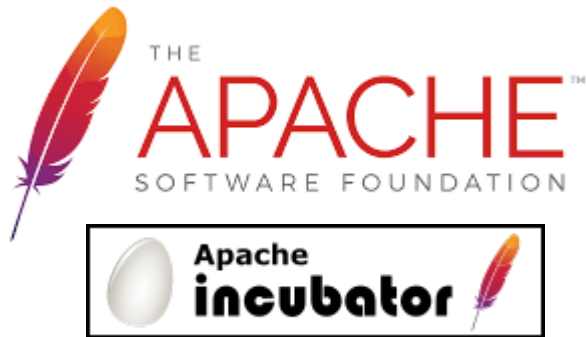
- Flexible
- Reliable
- High Performant
- Scalable

...*OLTP framework* that allows **BigData** apps to execute **ACID transactions** on top of **HBase**

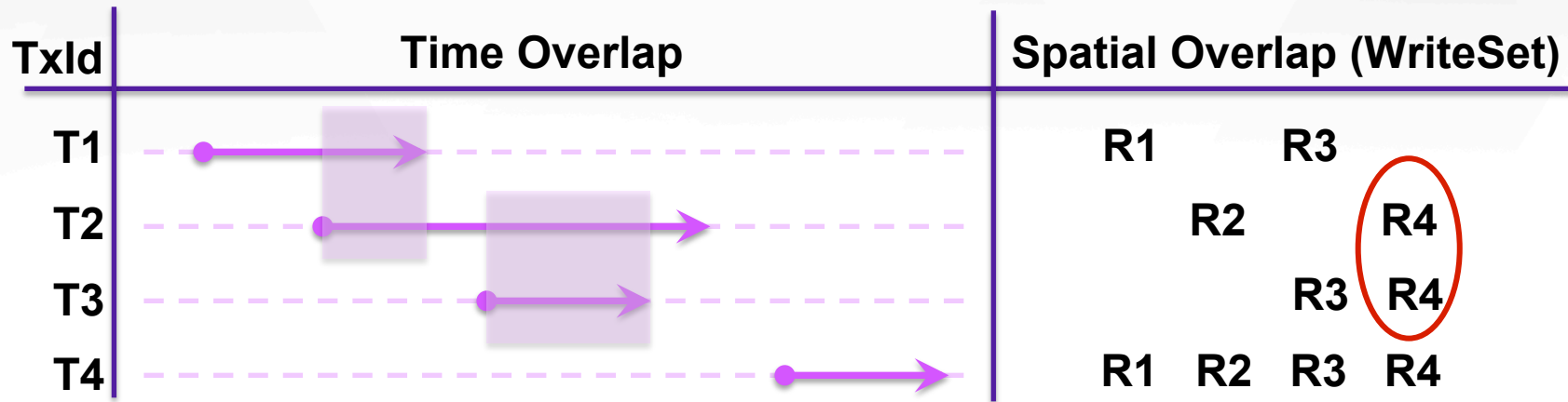
 +  = **Consistency in BigData Apps**

# Why use Omid?

- **Simplifies development** of apps requiring consistency
  - Multi-row/multi-table transactions on HBase
  - Simple & well-known interface
- **Good performance & reliability**
- **Lock-free**
- **Snapshot Isolation**
- **HBase is a blackbox**
  - No HBase code modification
  - No changes on table schemas
- **Used successfully at Yahoo**

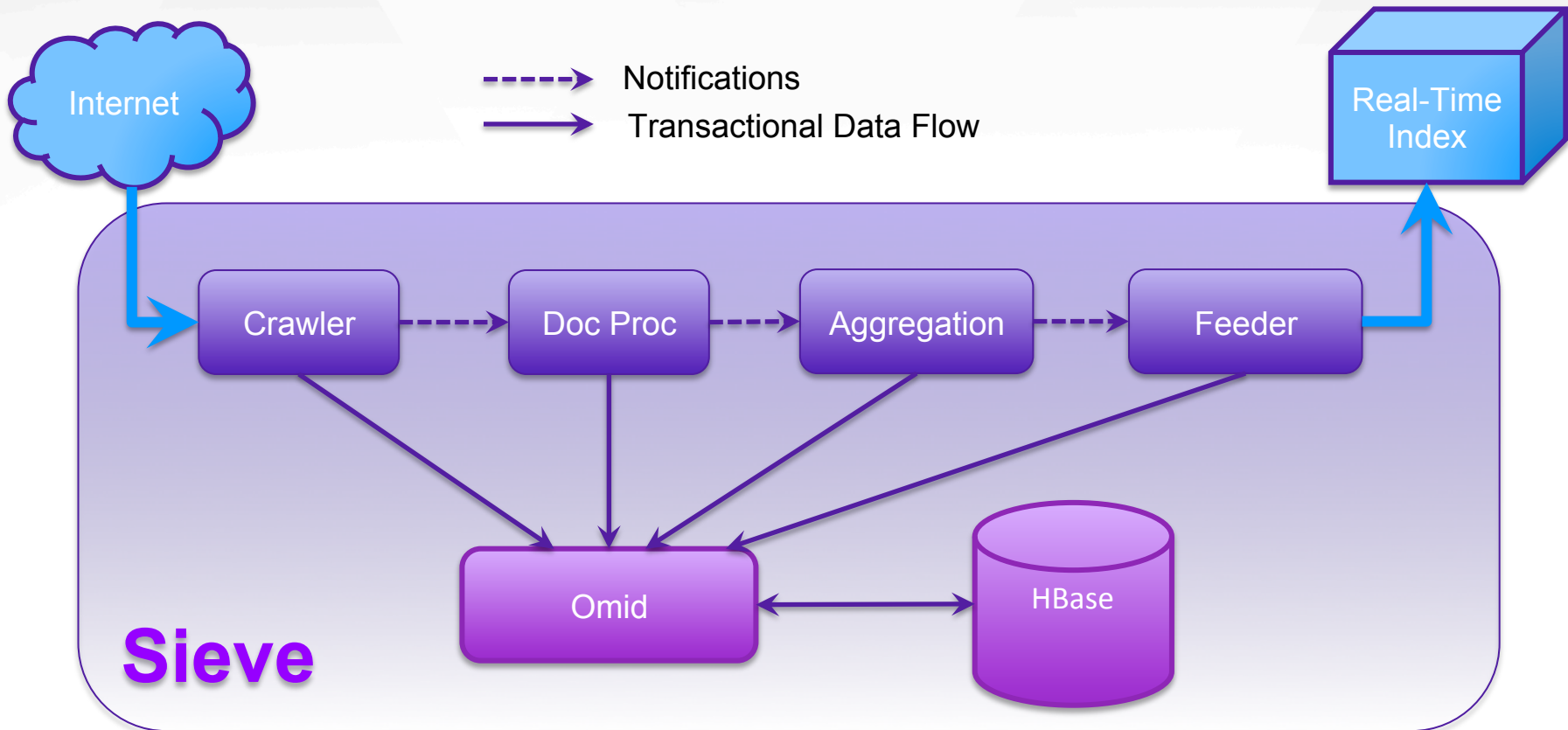


# Snapshot Isolation



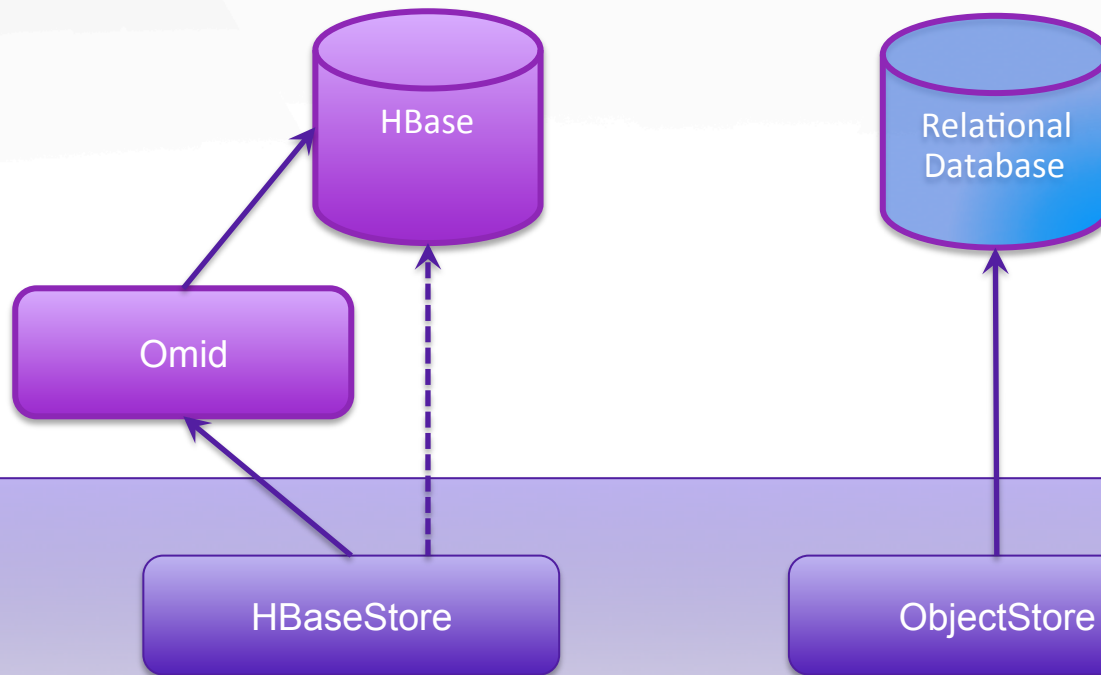
- Transaction T2 overlaps in time with T1 & T3, but spatially:
  - $T1 \cap T2 = \emptyset$
  - $T2 \cap T3 = \{ R4 \}$  Transactions T2 and T3 conflict
- Transaction T4 does not have conflicts

# Use Cases: Sieve @ Yahoo



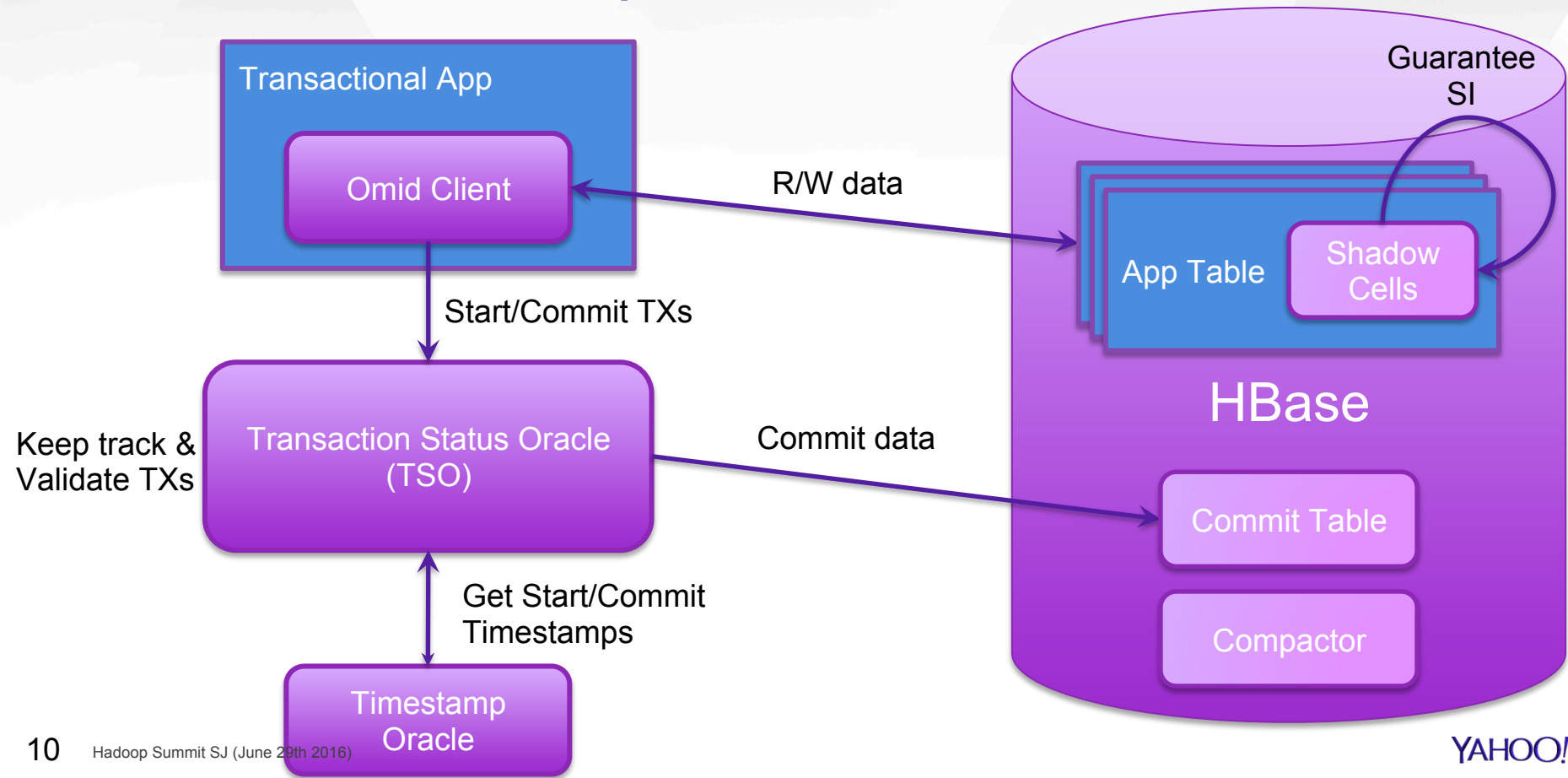


# Use Cases:



## Hive Metastore Thrift Server

# Architectural Components



# Client APIs

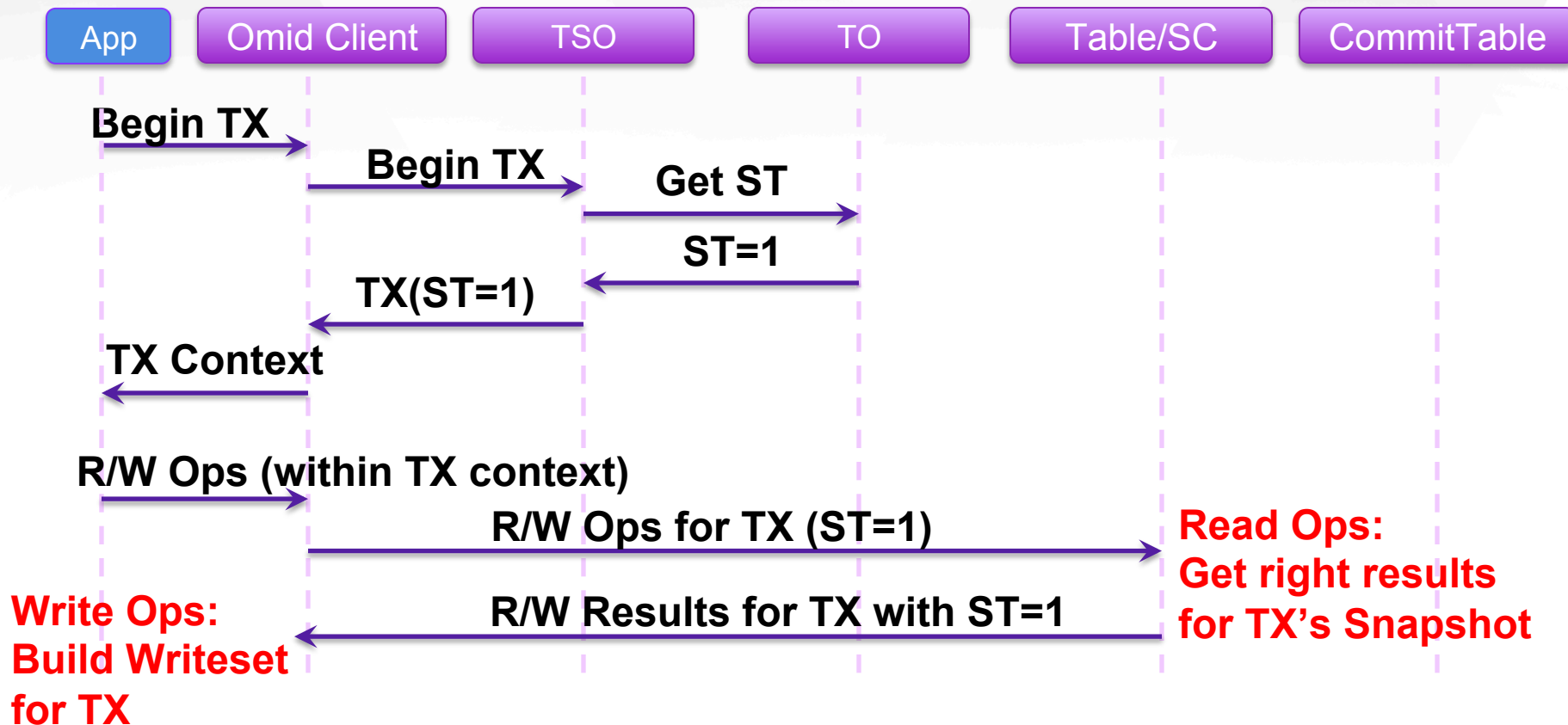
- Transaction Manager → Create Transactional contexts

```
Transaction begin();  
void commit(Transaction tx);  
void rollback(Transaction tx);
```

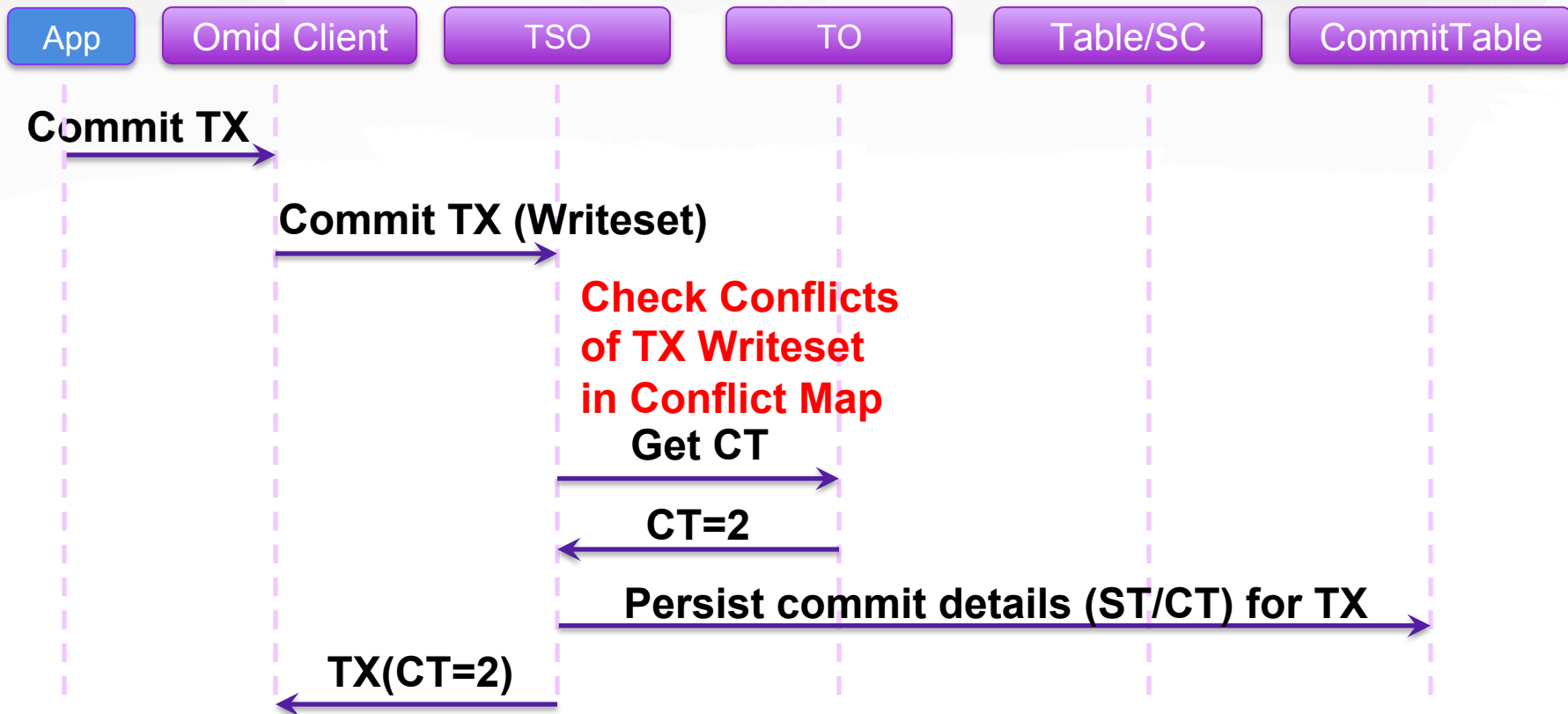
- Transactional Tables (TTable) → Data access

```
Result get(Transaction tx, Get g);  
void put(Transaction tx, Put p);  
ResultScanner getScanner(Transaction tx, Scan s);
```

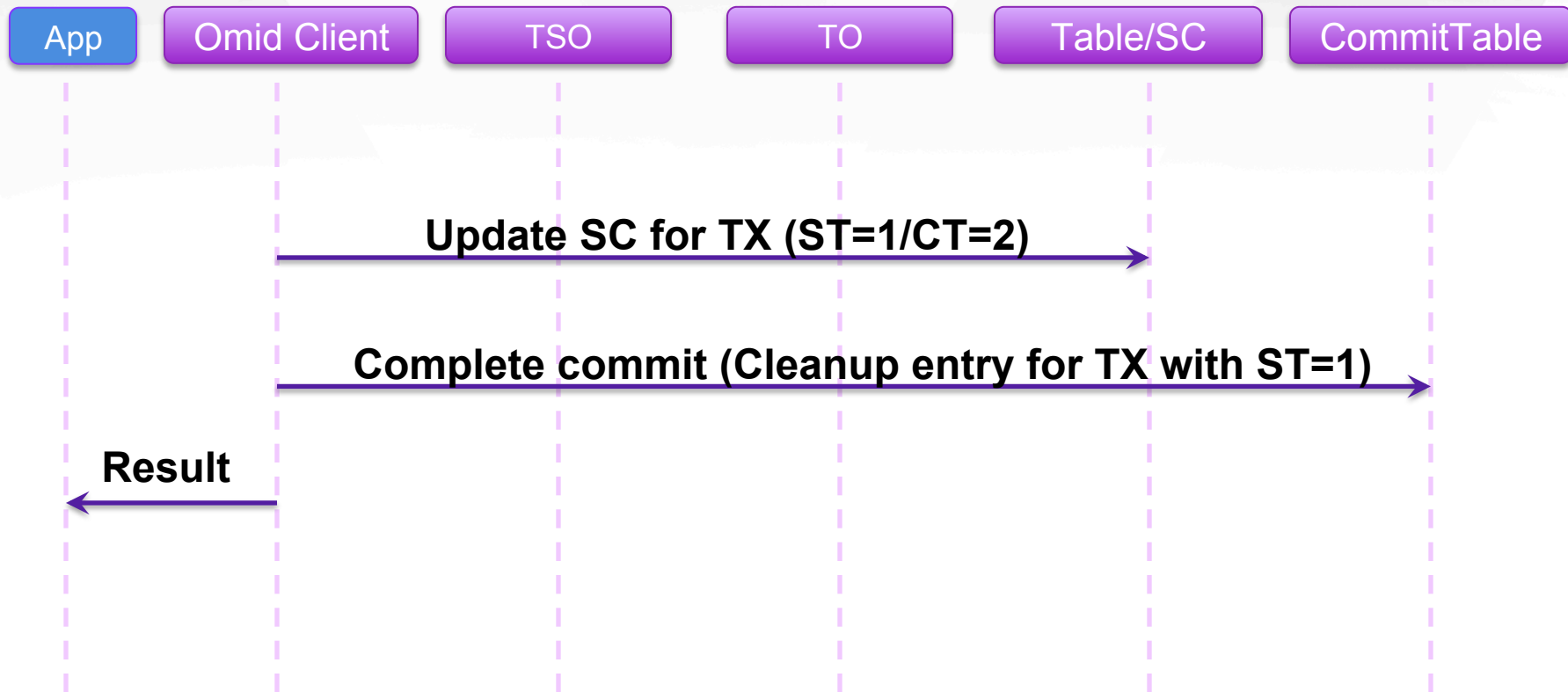
# TX Management (Begin TX phase)



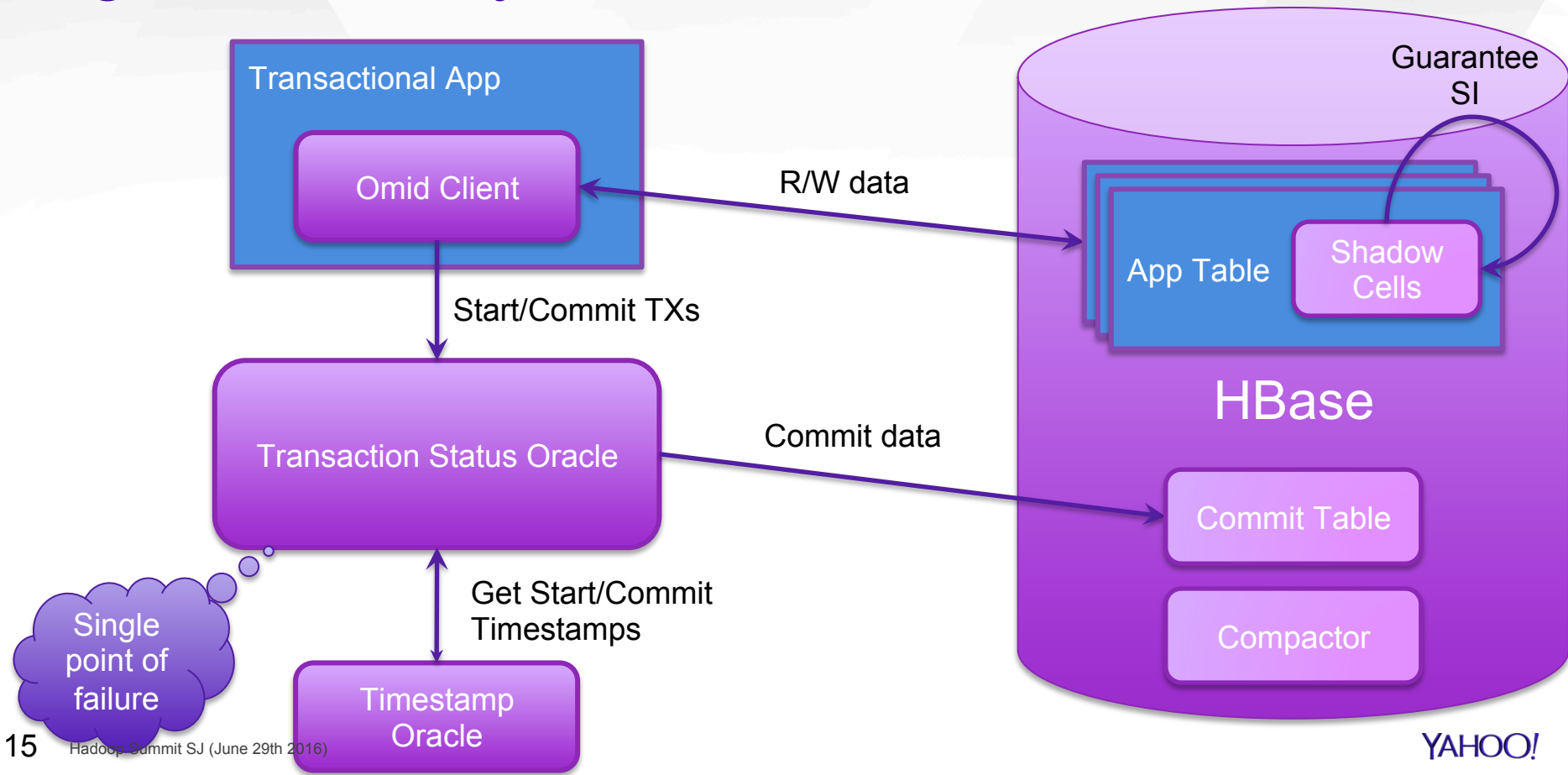
# TX Management (Commit TX Phase)



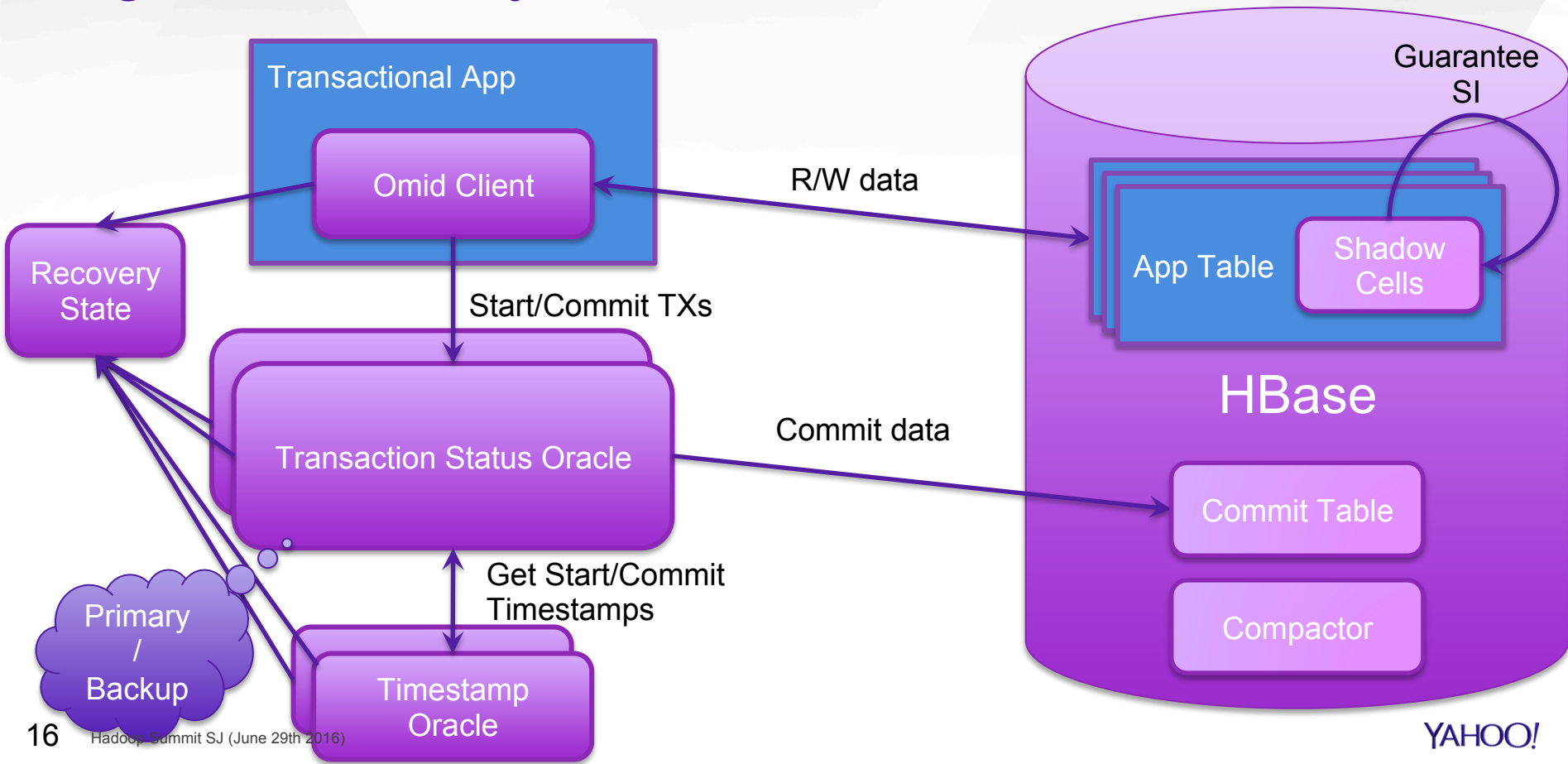
# TX Management (Complete TX Phase)



# High Availability

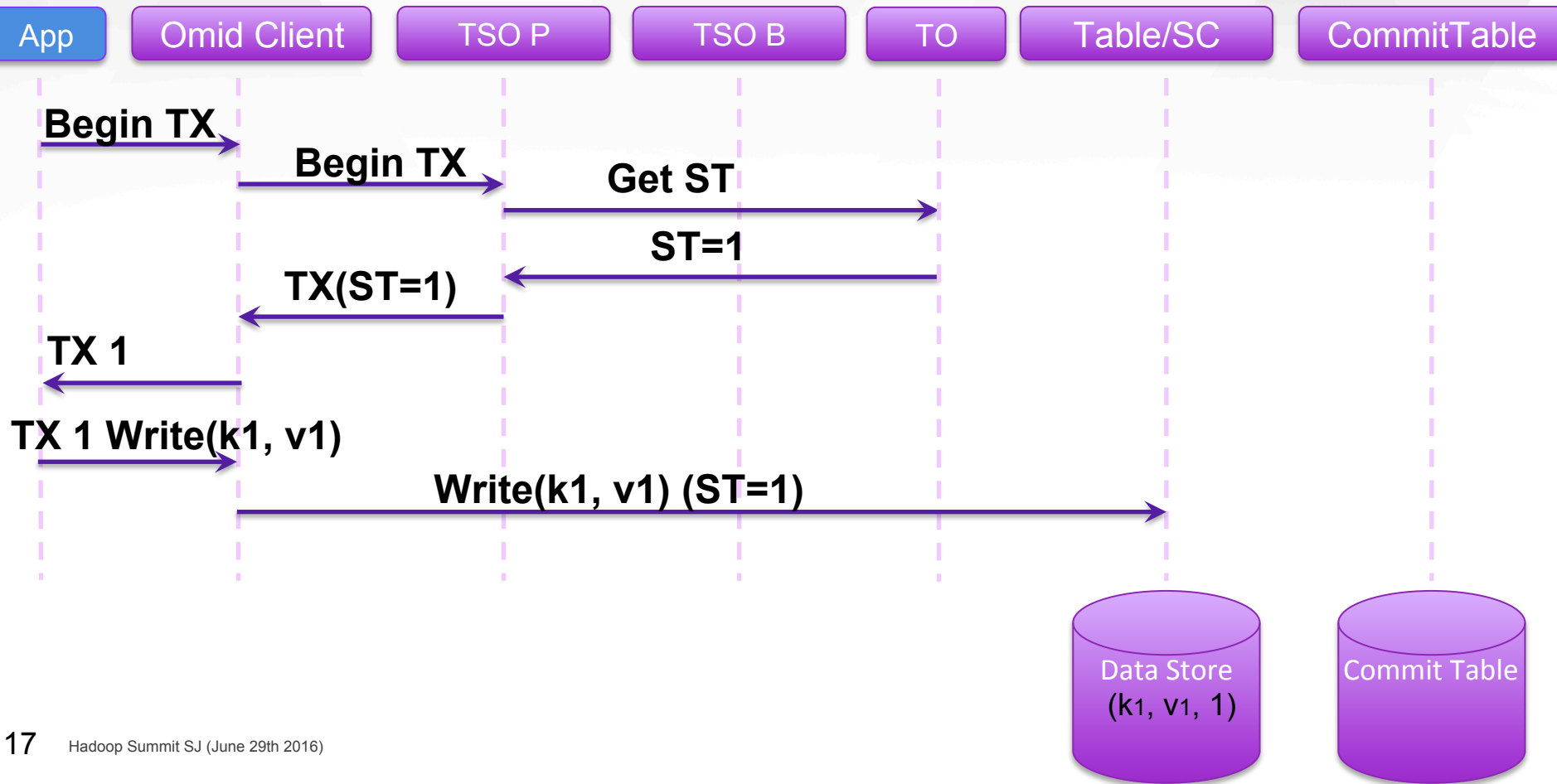


# High Availability

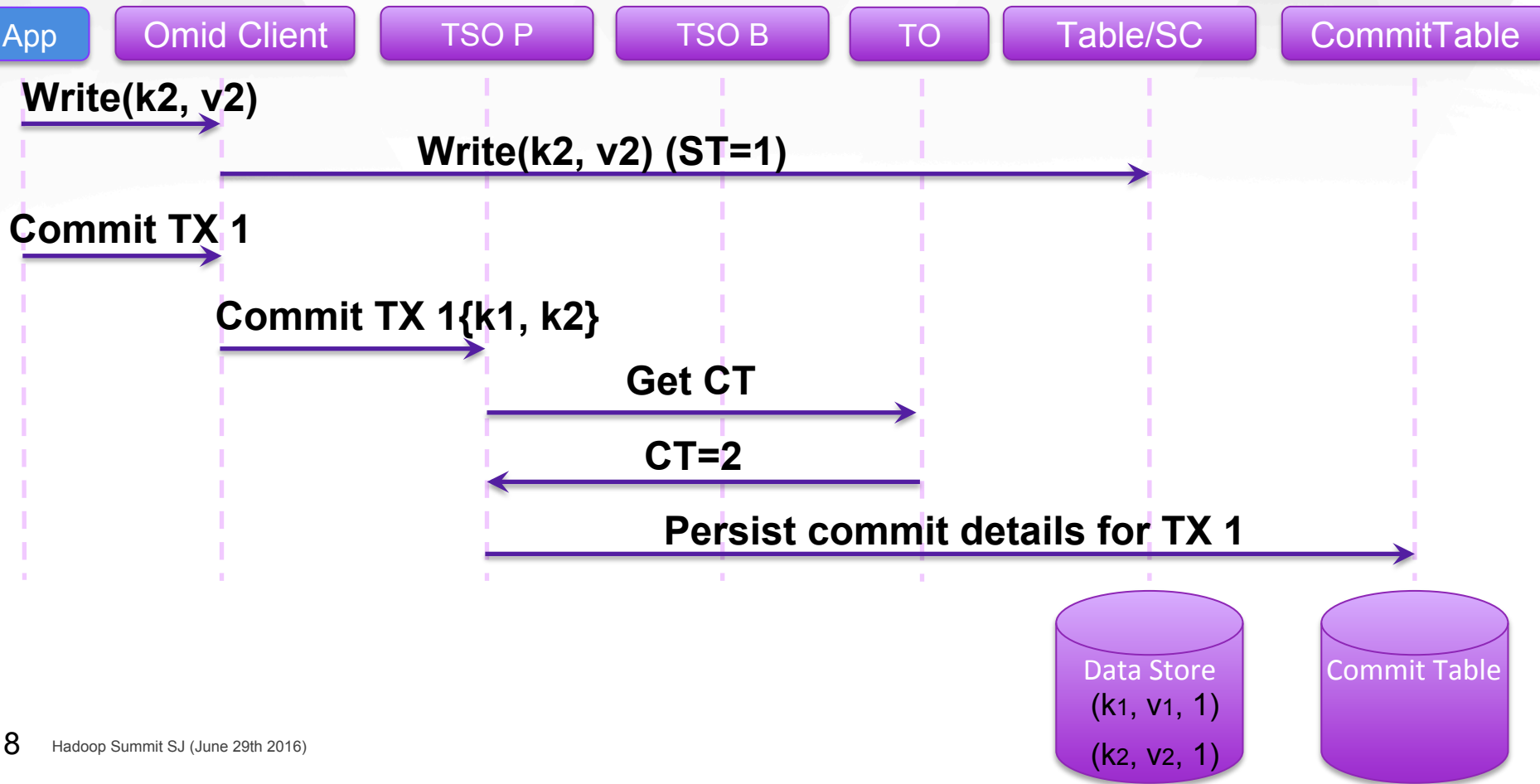




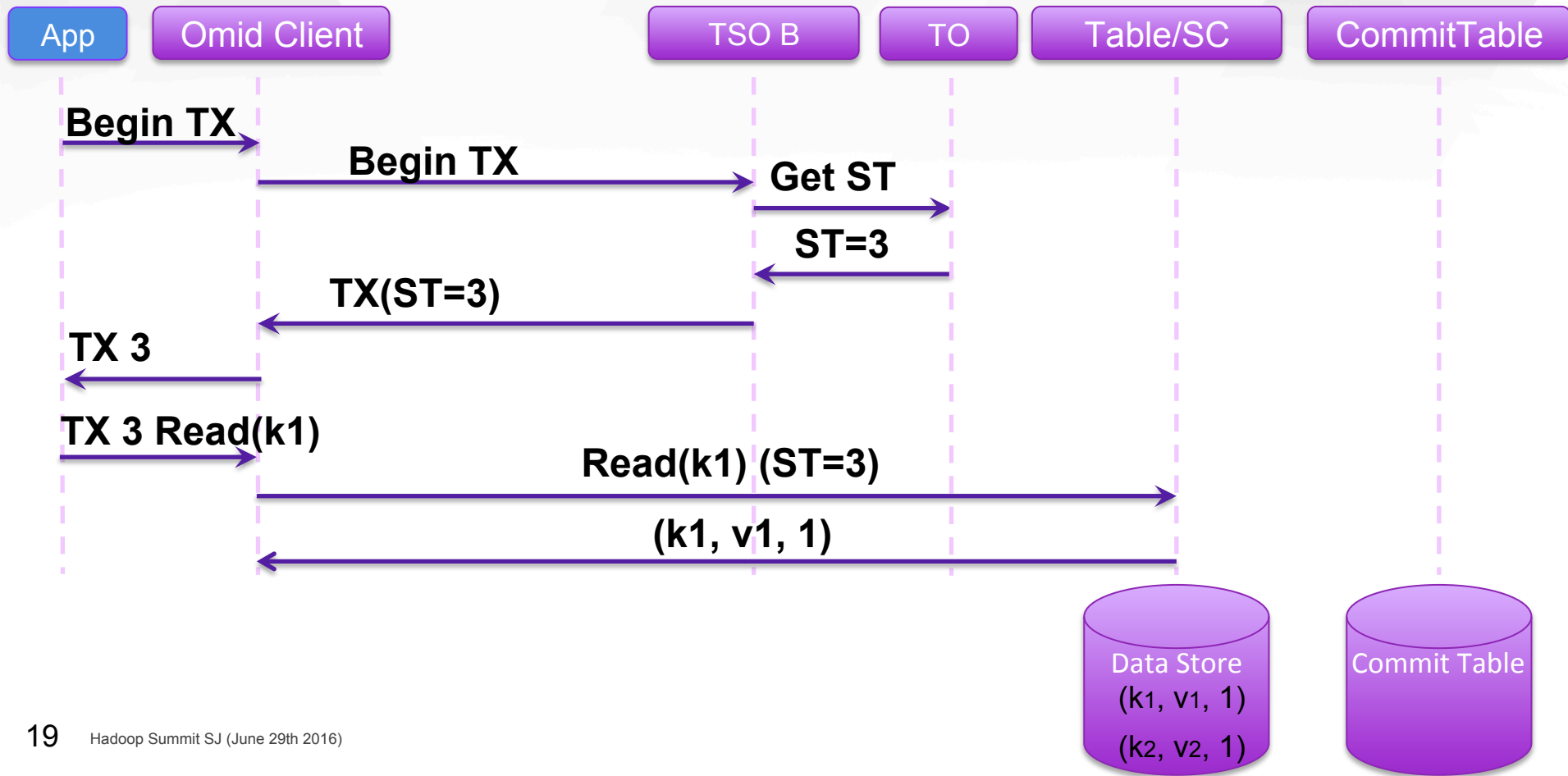
# High Availability – Failing Scenario



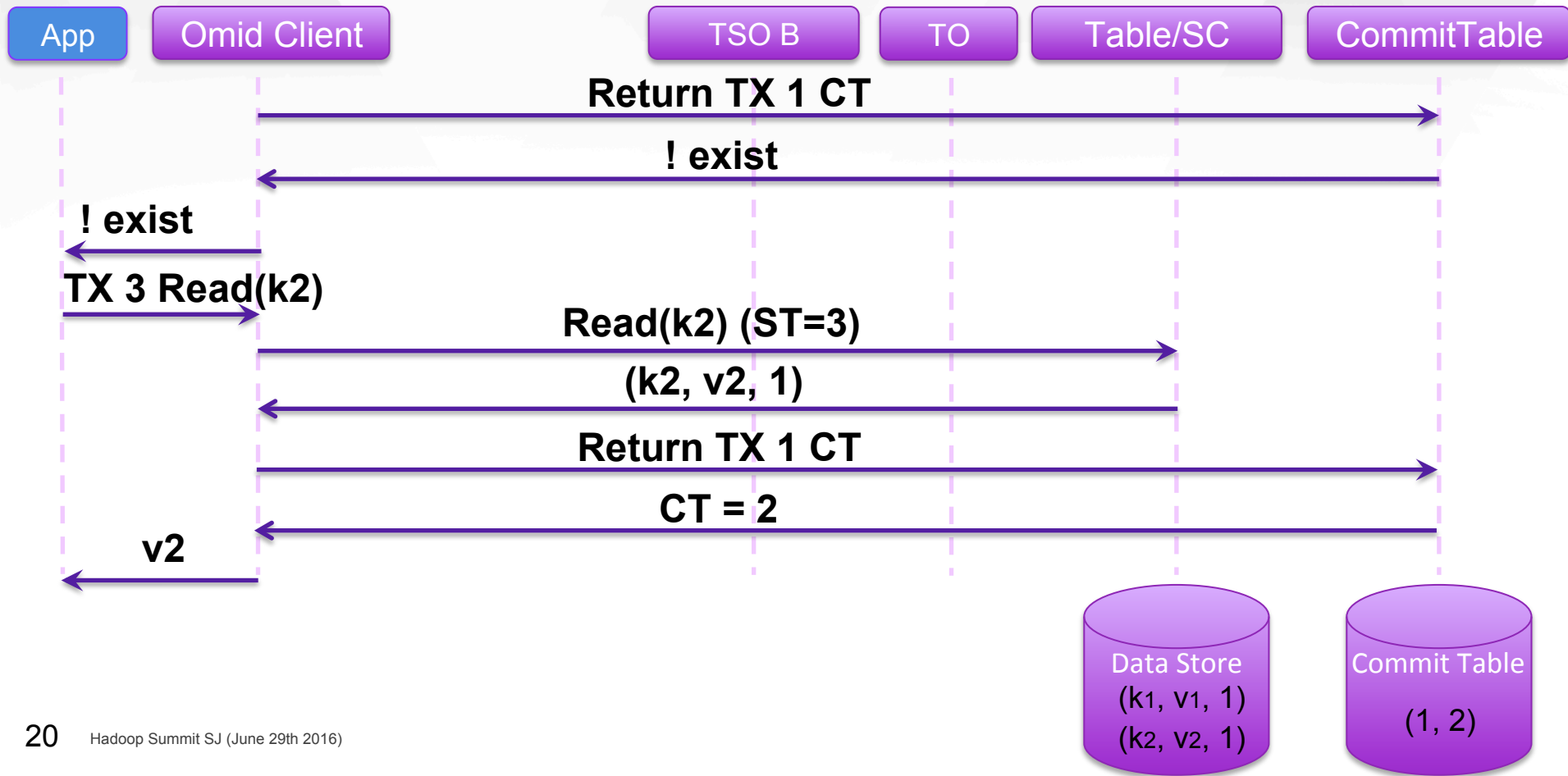
# High Availability – Failing Scenario



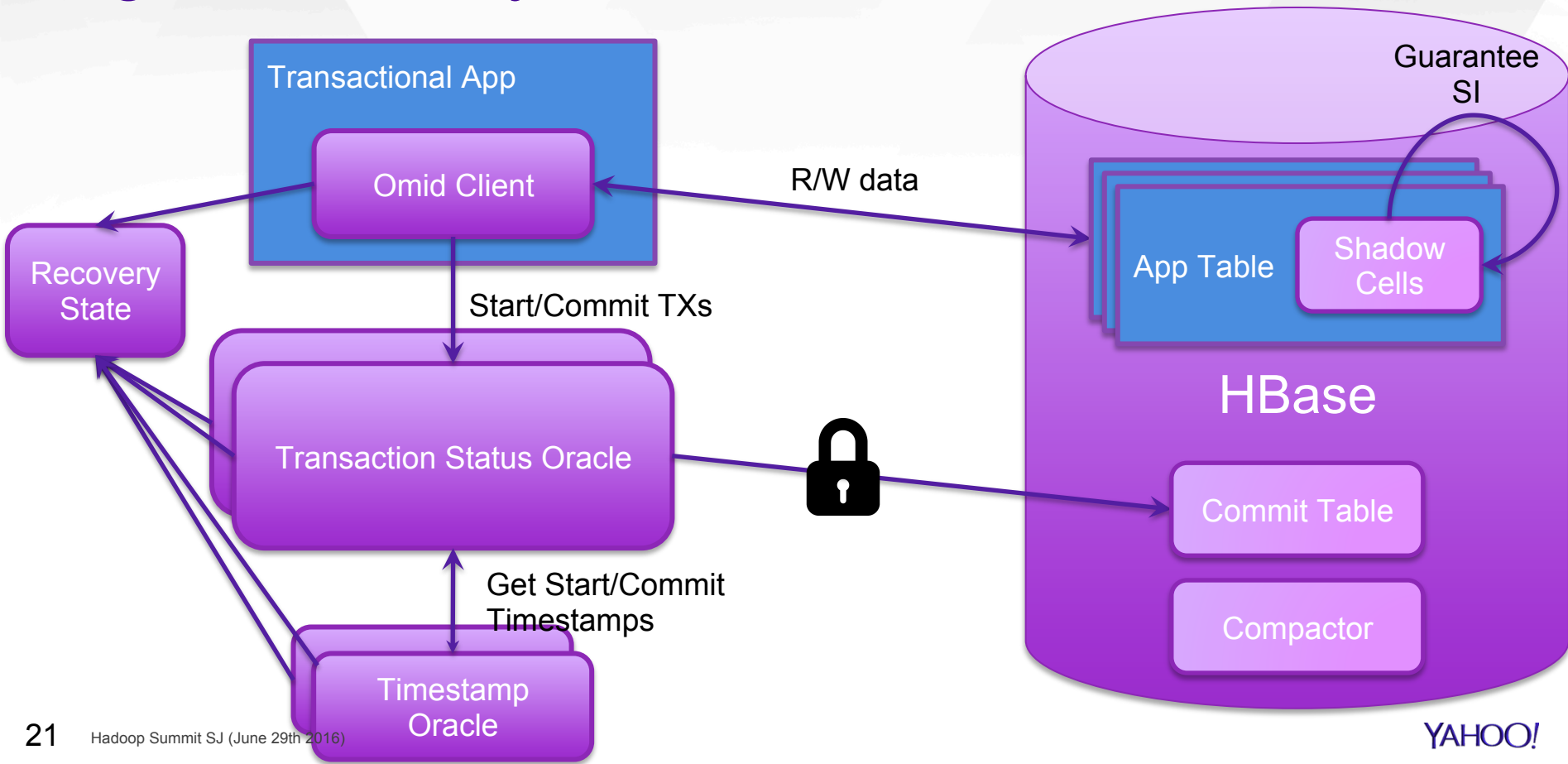
# High Availability – Failing Scenario



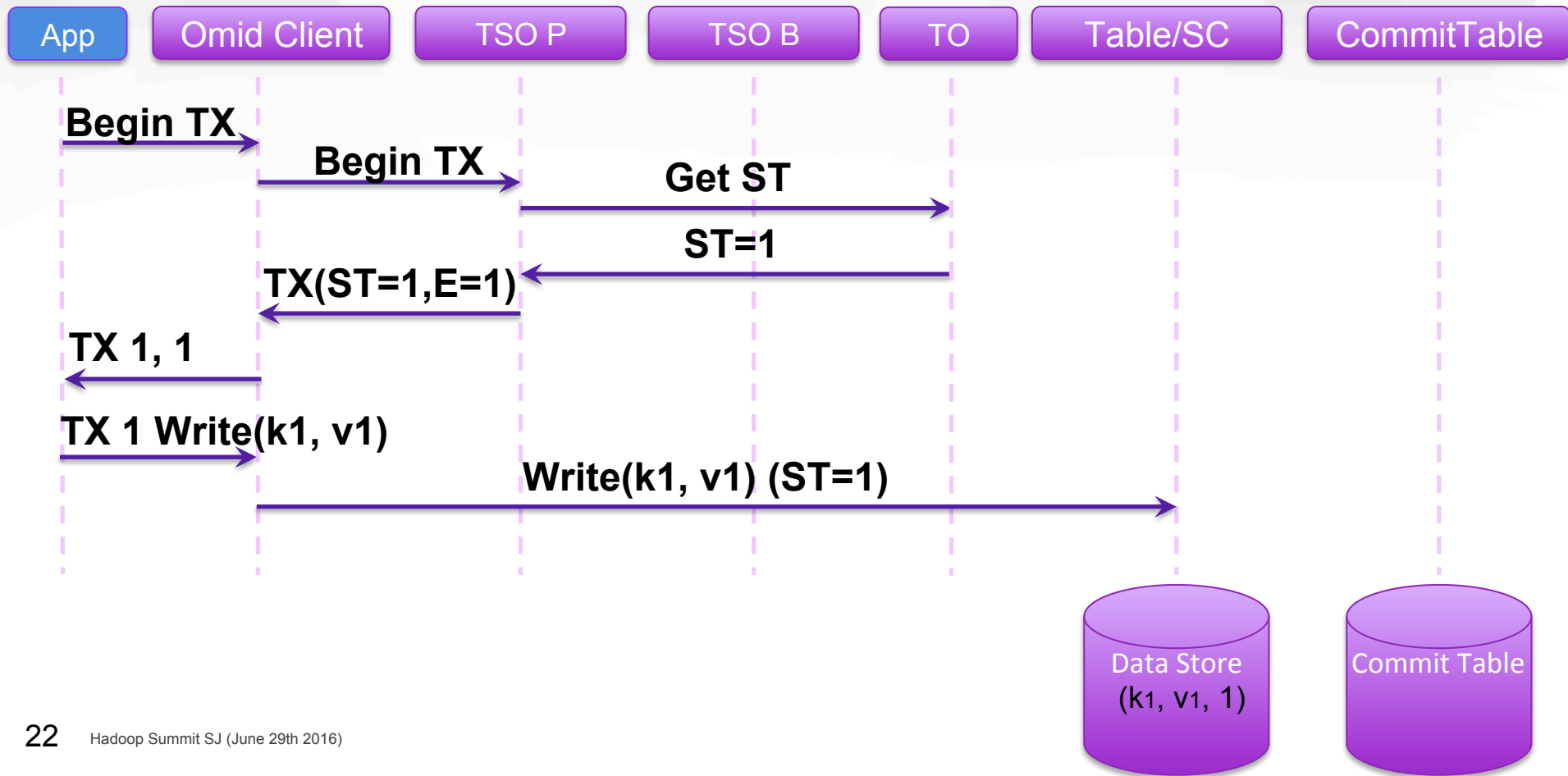
# High Availability – Failing Scenario



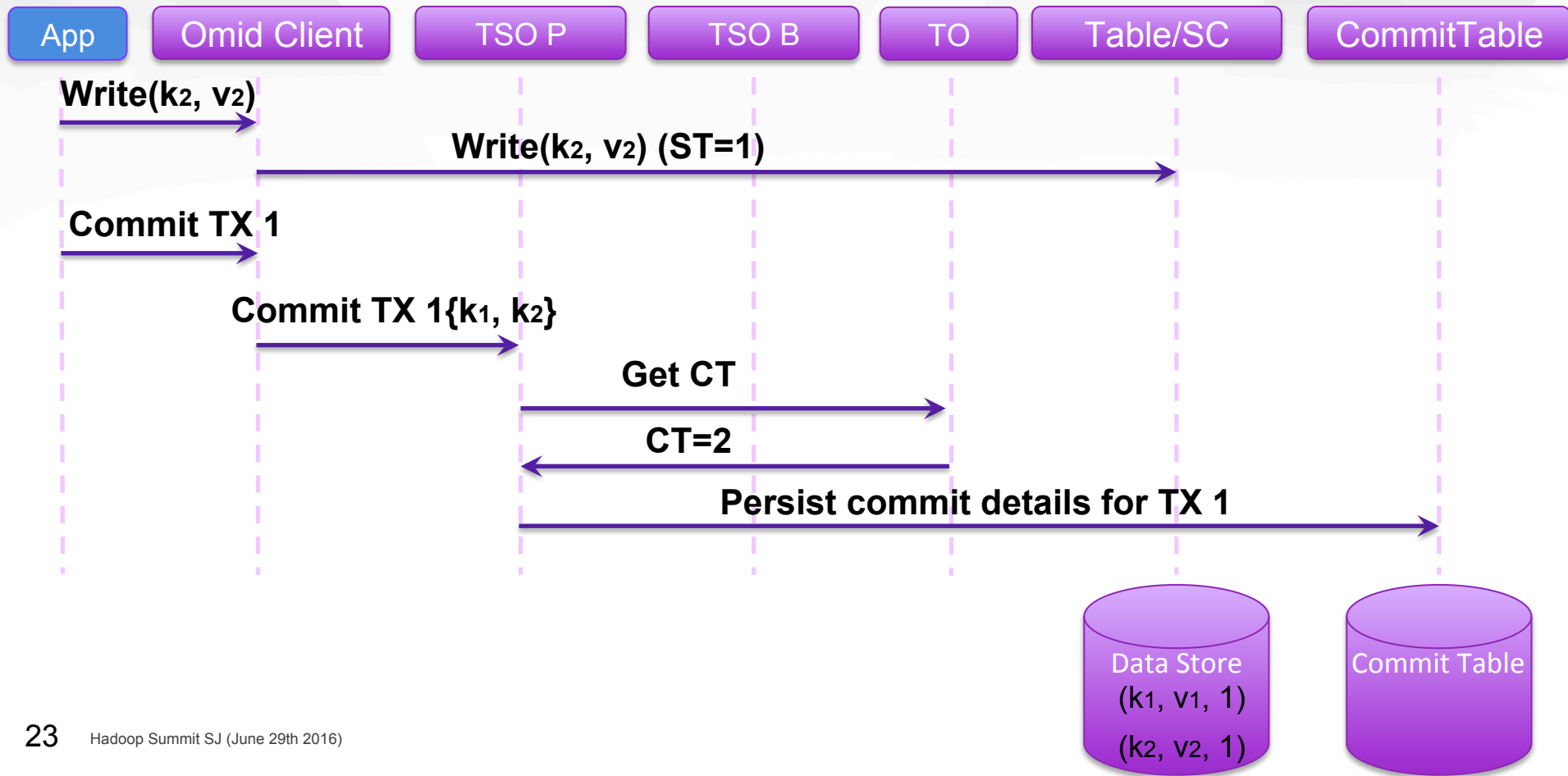
# High Availability



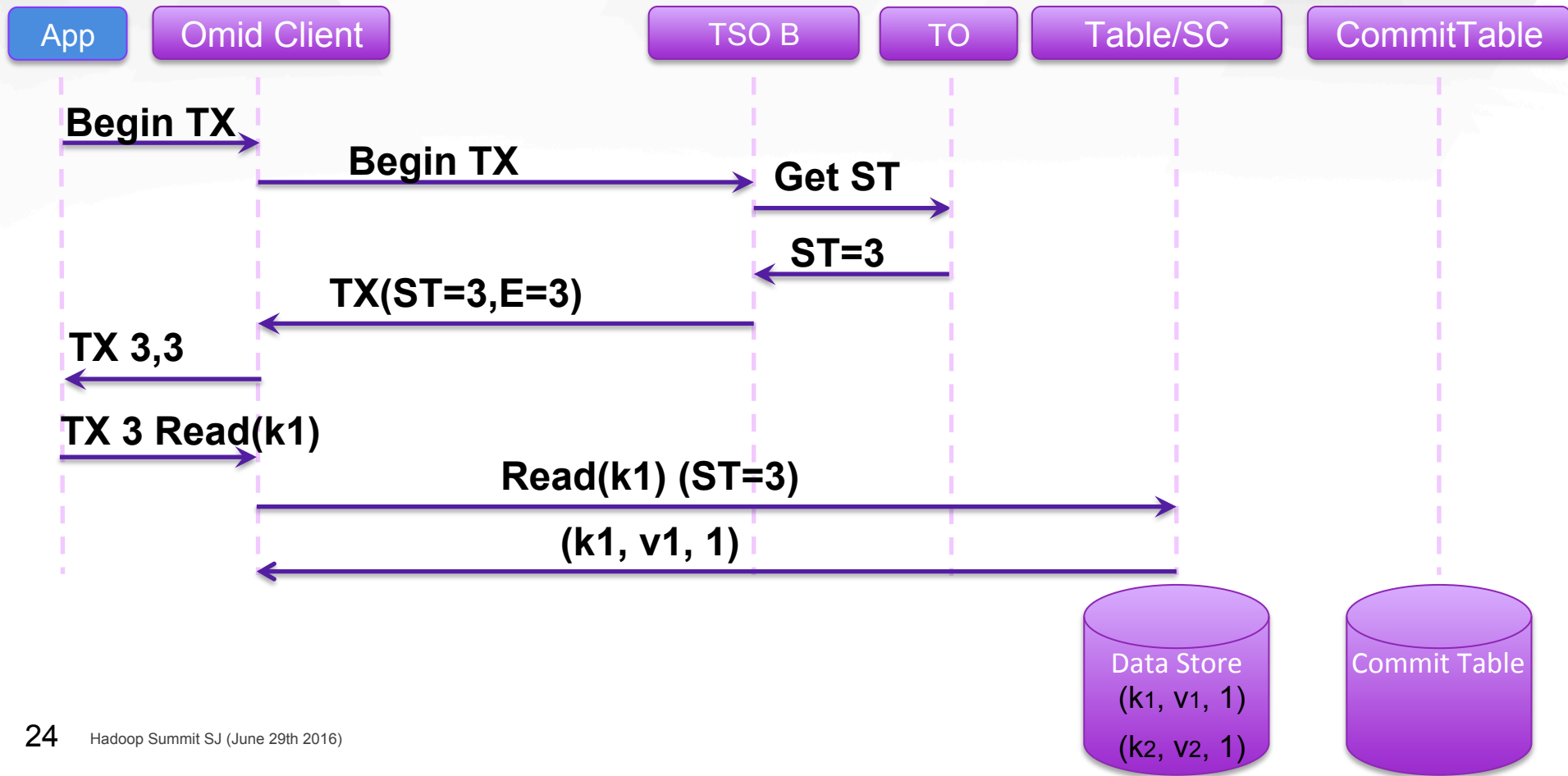
# High Availability – Solution



# High Availability – Solution

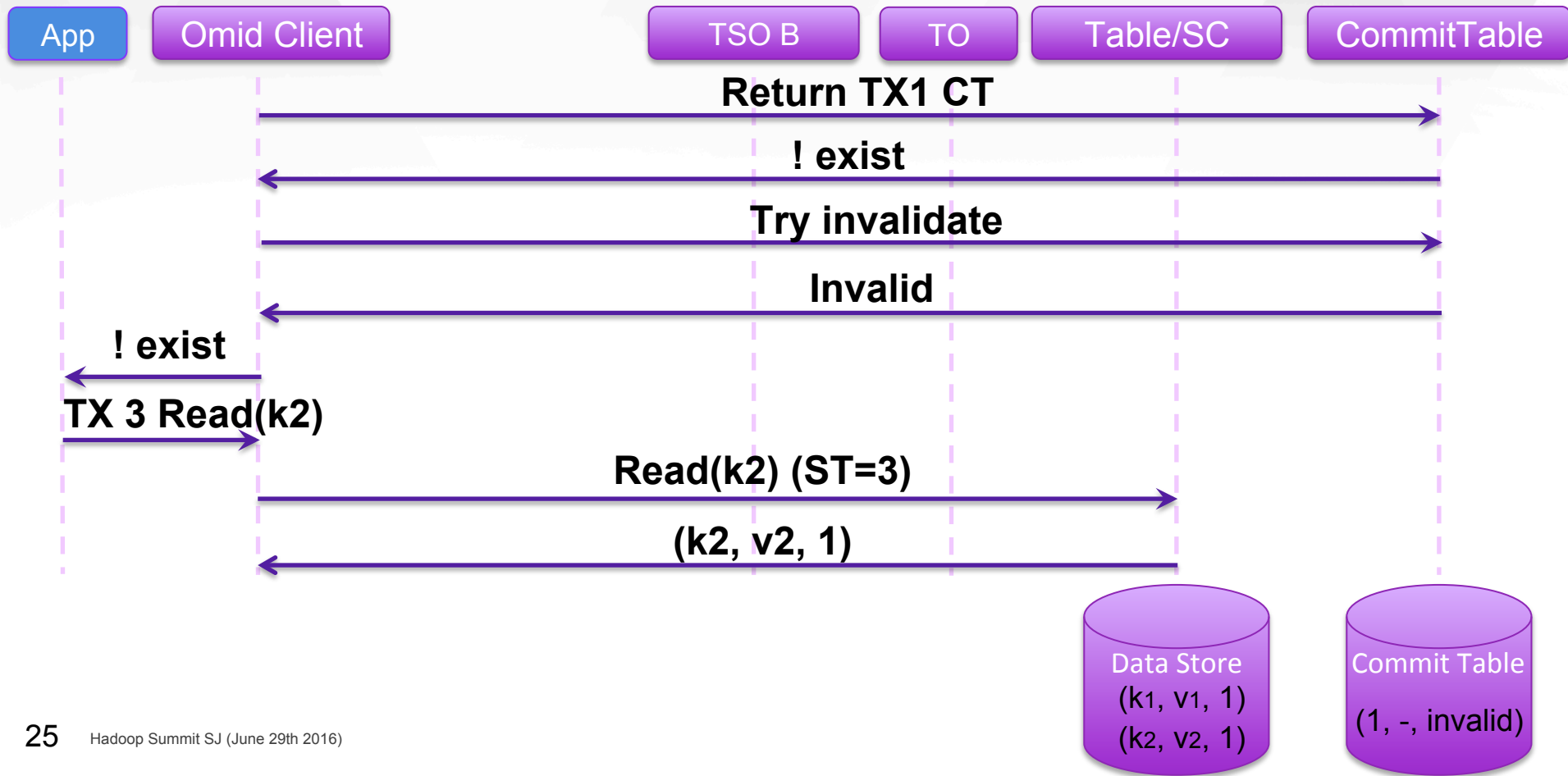


# High Availability – Solution

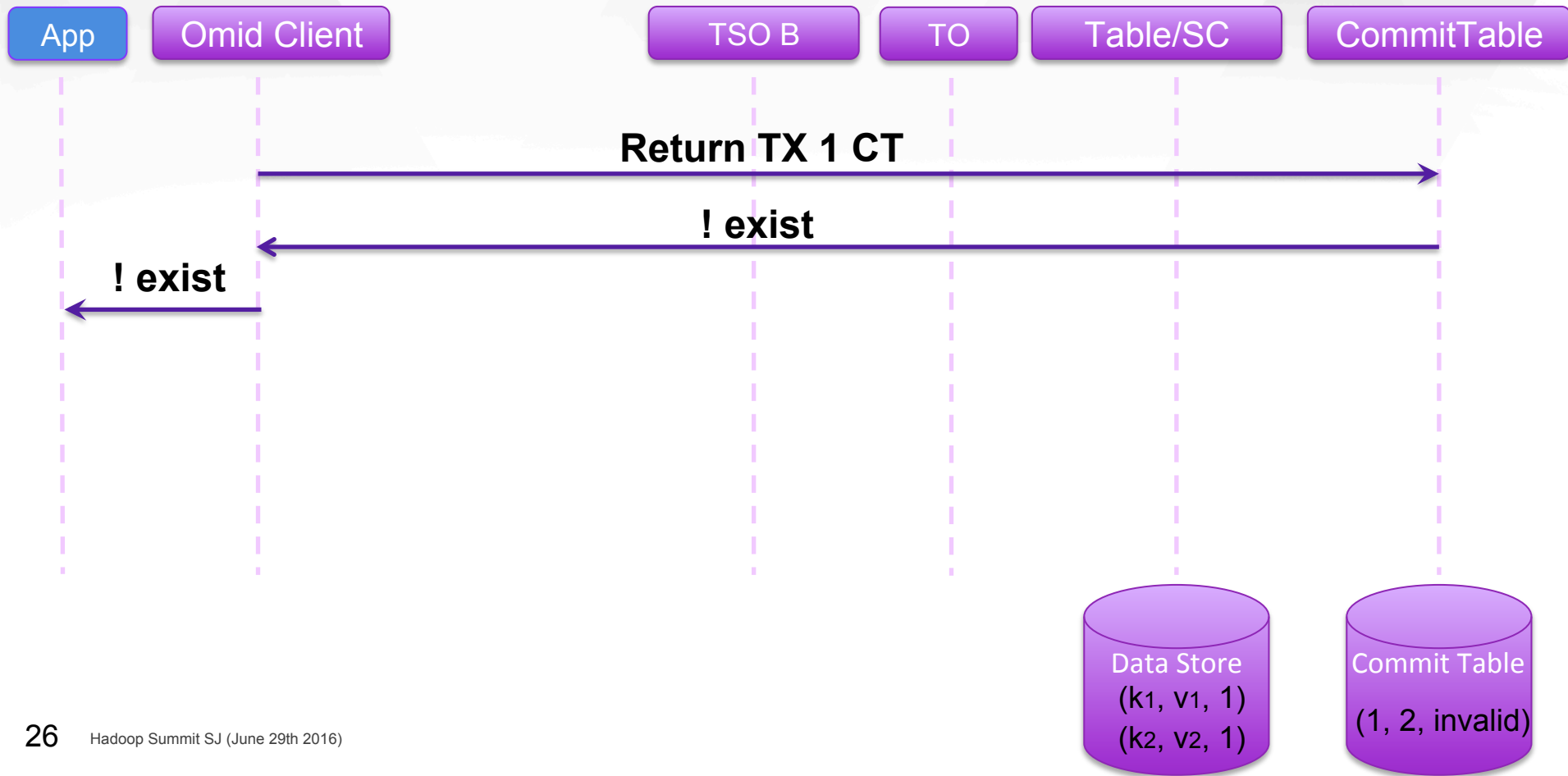




# High Availability – Solution



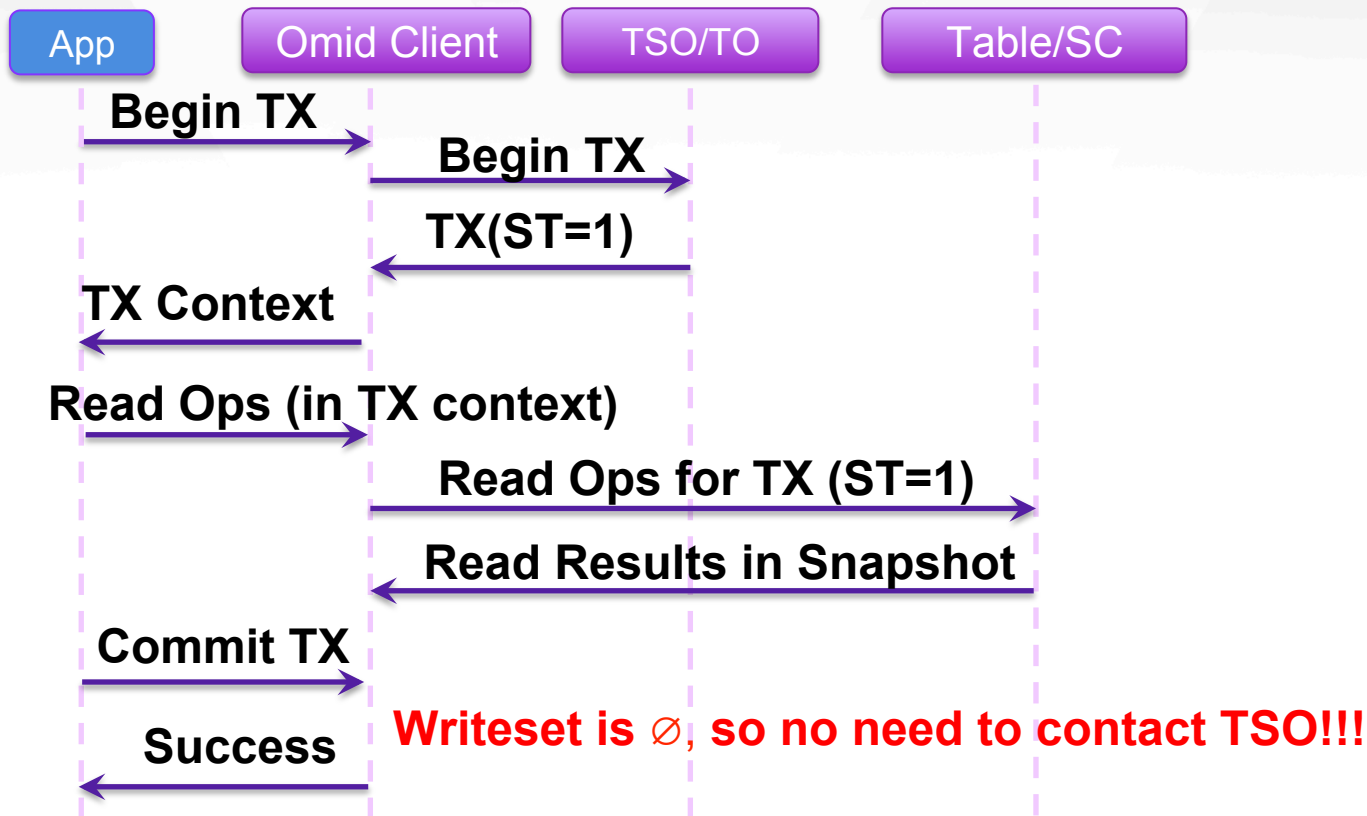
# High Availability – Solution



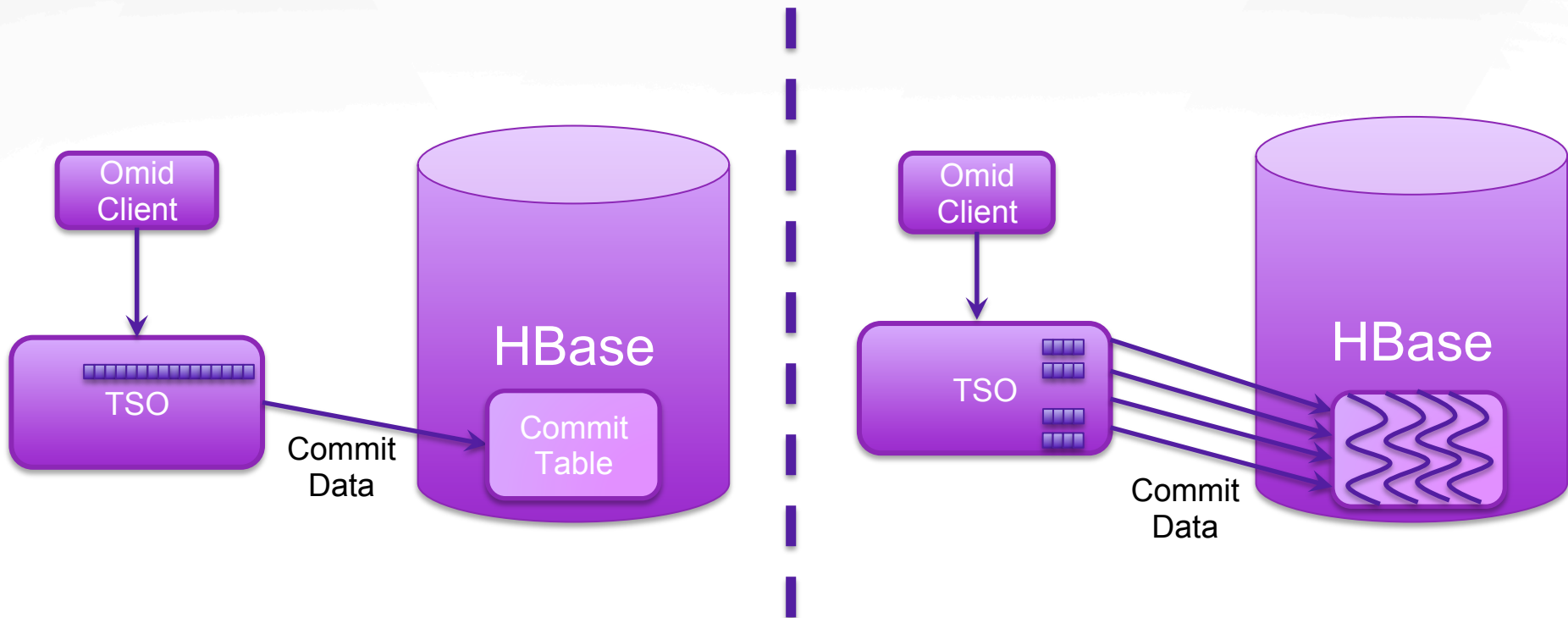
# High Availability

- **No runtime overhead** in mainstream execution
  - Minor overhead after failover
- TSO uses **regular writes**
- **Leases** for leader election
  - Lease status check before/after writing to Commit Table

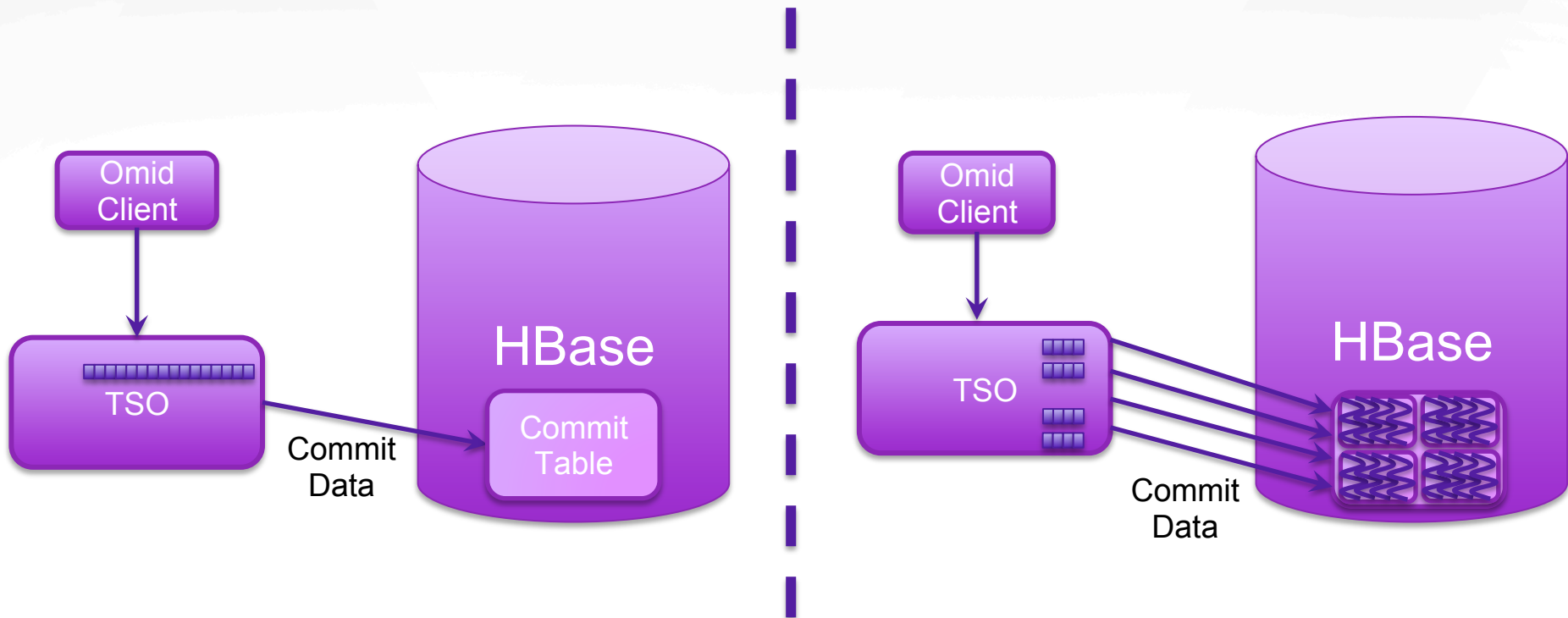
# Perf. Improvements: Read-Only Txs



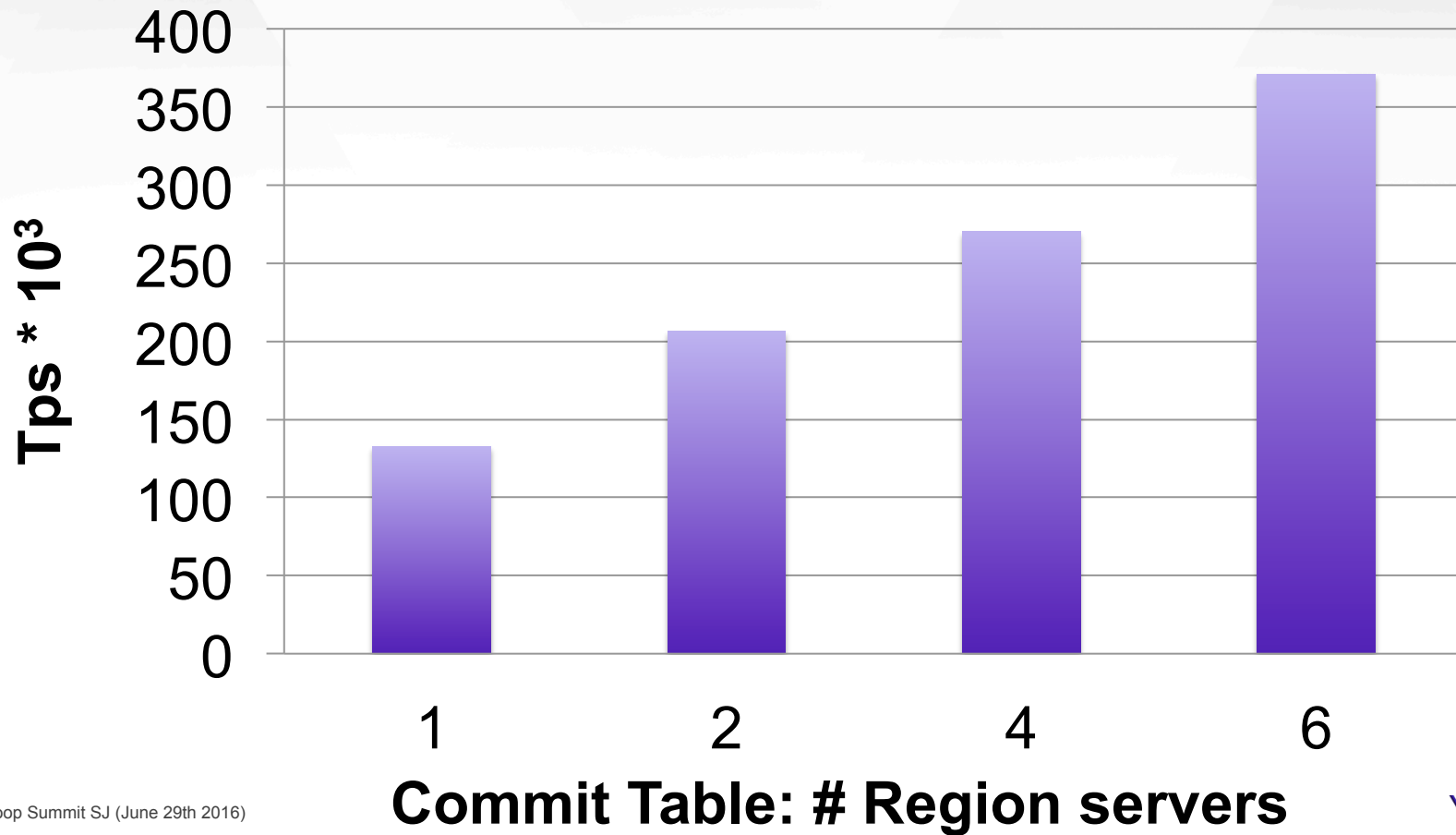
# Perf. Improvements: Commit Table Writes



# Perf. Improvements: Commit Table Writes



# Omid Throughput with Improvements



# Summary



## ■ Transactions in NoSQL

- Use cases in incremental big data processing
- *Snapshot Isolation*: Scalable consistency model

## ■ Omid

- Web-scale TPS for HBase
- Reliable and performant
- Battle-tested



<http://omid.incubator.apache.org/>



# Questions?



[@ApacheOmid](#)  
[Apache Omid Incubator Page](#)